Problems

- 1. Solve each of the following problems. Note: correctly solving these problems counts for passing LO1.
 - a. Evaluate $5^{40} \mod 17$ without the help of a calculator. (10 pts)
 - b. In the Strassen-Solovay test, is 8 and witness or accomplice for n = 15? Show work in computing both the left and right sides of the mod-15 congruence. (15 pts)
- 2. Solve each of the following problems. Note: correctly solving these problems counts for passing LO2.
 - a. Use the Master Theorem to determine the growth of T(n) if it satisfies the recurrence $T(n) = 4T(n/2) + n^{\log_3 10} \log^2 n$. (10 pts)
 - b. Use the substitution method to prove that, if T(n) satisfies

$$T(n) = 4T(n/2) + n^2,$$

Then $T(n) = O(n^2 \log n)$. (15 pts)

- 3. Solve each of the following problems. Note: correctly solving these problems counts for passing LO3.
 - a. Consider the following algorithm called **multiply** for multiplying two *n*-bit binary numbers x and y. In what follows, we assume n is even. Let x_L and x_R be the leftmost n/2 and rightmost n/2 bits of x respectively. Define y_L and y_R similarly. Let P_1 be the result of calling **multiply** on inputs x_L and y_L , P_2 be the result of calling **multiply** on inputs x_R and y_L , P_2 be the result of calling **multiply** on inputs x_R and y_R , and P_3 the result of calling **multiply** on inputs $x_L + x_R$ and $y_L + y_R$. Then return the value $P_1 \times 2^n + (P_3 P_1 P_2) \times 2^{n/2} + P_2$. Prove that the returned value does in fact equal xy. (15 pts)
 - b. Use Strassen's products $P_1 = a(f h) = af ah$, $P_2 = (a + b)h = ah + bh$, $P_3 = (c + d)e = ce + de$, $P_4 = d(g e) = dg de$, $P_5 = (a + d)(e + h) = ae + ah + de + dh$, $P_6 = (b d)(g + h) = bg + bh dg dh$, $P_7 = (a c)(e + f) = ae ce cf + af$. to compute the matrix product

$$\left(\begin{array}{rrr}1 & -3\\-4 & 5\end{array}\right)\left(\begin{array}{rrr}3 & -1\\2 & 4\end{array}\right)$$

Show all work. (10 pts)

4. Recall that, for integers a, b, and c, $(a, b) \mid c$ iff there exist integer constants x and y for which

$$ax + by = c.$$

Use this fact to prove the following.

a. If the equation

 $ax \equiv b \mod m$

has a solution then $(a, m) \mid b$. (12 pts)

b. If $(a, m) \mid b$, then the equation

 $ax \equiv b \mod m$

has a solution. (13 pts)

- 5. Show how to multiply the complex numbers a + bi and c + di using only three multiplications of real numbers. The algorithm should take a, b, c, and d as input, and produce the real component ac - bd and imaginary component ad + bc. Note that the straightforward approach requires four multiplications. We seek a more clever approach. (25 pts)
- 6. Given an array a of n positive integers, the maximum window area (MWA) of a is defined as the maximum of

$$(j-i+1) \min_{i \leq k \leq j} (a[k]),$$

taken over all combinations i and j for which $0 \le i \le j \le n-1$. For example if a = 3, 3, 1, 7, 4, 2, 4, 6, 1, then MWA(a) = 10 via i = 3 and j = 7, since the minimum value in this range is a[5] = 2, and (7 - 3 + 1)(2) = 10. One algorithm for finding MWA(a) is to consider all n^2 possible combinations of i and j and keep track of the combination that produces the maximum window area. But this algorithm has quadratic running time.

- a. Describe a divide-and-conquer algorithm that achieves an improved running time. Clearly define the steps of your algorithm. Use any results from the Recurrences lecture to analyze its running time. (15 pts)
- b. Demonstrate your algorithm on the array a = 2, 6, 3, 7, 5, 4, 6, 2, 1, 7 for the top level of recursion only. For example, if your algorithm makes two recursive calls, then (without working through the algorithm) provide their solutions and show the steps of the combine portion of the algorithm working at the top level. (10 pts)