## CECS 329, LO Pre-Final Assessment, Dr. Ebert

NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION
ALLOWED. Submit solutions to AT MOST 6 PROBLEMS. Please use BOTH SIDES of each answer sheet to save paper. Make sure your name and Class ID are on each answer sheet.

LO1. Answer the following.
(a) Provide the definition of what it means to be a mapping reduction from decision problem $A$ to decision problem $B$.
(b) The simple graph $G=(V, E)$ shown below is an instance of the Clique problem with $k=4$. Draw $f(G, k)$, where $f$ is the mapping reduction from Clique to Half Clique provided in the Mapping Reducibility lecture exercises.

(c) Verify that both $G$ and $f(G, k)$ are either both positive or both negative instances of their respective problems. Justify your answer.

LO2. Recall that the 3-Dimensional Matching (3DM) decision problem takes as input three sets $A$, $B$, and $C$, each having size $n$, along with a set $S$ of triples of the form ( $a, b, c$ ) where $a \in A$, $b \in B$, and $c \in C$. We assume that $|S|=m \geq n$. The problem is to decide if there exists a subset $T \subseteq S$ of $n$ triples for which each member from $A \cup B \cup C$ belongs to exactly one of the triples. Complete the following steps to show that 3DM is a member of NP.
(a) Define a certificate for an instance $(A, B, C, S)$ of 3DM. Note: letter $C$ is already being used, so use a different letter to represent the certificate.
(b) Provide a semi-formal verifier algorithm that takes as input 3DM instance ( $A, B, C, S$ ) and a certificate and returns 1 iff the certificate is valid.
(c) Provide size parameters for 3DM.
(d) Provide the verifier's running time and defend your answer.

LO3. An instance $\mathcal{C}$ of 3SAT consists of clauses $c_{1}=\left(x_{1}, \bar{x}_{2}, x_{3}\right), c_{2}=\left(\bar{x}_{2}, x_{3}, x_{4}\right), c_{3}=\left(\bar{x}_{1}, x_{2}, \bar{x}_{4}\right)$, and $c_{4}=\left(\bar{x}_{1}, \bar{x}_{3}, x_{4}\right)$. Answer the following questions about the mapping reduction $f(\mathcal{C})=(G, a, b)$ provided in lecture from 3SAT to Directed Hamilton Path and applied to instance $\mathcal{C}$.
(a) Consider the vertices $l c_{2}$ and $r c_{2}$ that lie in the chain that connects the left and right halves of the $x_{2}$-diamond of $G$. Describe the orientation of the edge between $l c_{2}$ and vertex $c_{2}$ as well as the orientation of the edge between $r c_{2}$ and $c_{2}$. Draw a figure with the three vertices and the two edges and explain why you chose the particular orientations.
(b) Which diamond has no edges connected to vertex $c_{2}$ ? Explain.
(c) Given that $\alpha=\left(x_{1}=x_{2}=0, x_{3}=1, x_{4}=0\right)$ satisfies $\mathcal{C}$, provide the direction (left or right) that the corresponding Hamilton Path will take in each of the four diamonds of $G$. Hint: moving left means the path will move across the diamond from left to right.

LO4. Answer the following.
(a) Write either P, NP, or co-NP next to each of the following decision problems in terms of which one best characterizes the complexity of the problem. Note: there is only one best answer for each. Two points each.
i. An instance of Independent $\operatorname{Set}$ is a graph $G=(V, E)$ and a nonnegative integer $k \geq 0$, and the problem is to decide if $G$ has an independent set of size $k$.
ii. An instance of Tautology is a Boolean formula $F$ and the problem is to decide if $F$ evaluates to 1 for every Boolean assignment of its variables.
iii. An instance of Bounded Clique Size is a graph $G=(V, E)$ and a nonnegative integer $k \geq 0$, and the problem is to decide if every clique of $G$ has a size that does not exceed $k$.
iv. An instance of Triangle is a graph $G=(V, E)$ and the problem is to decide if $G$ has 3-Clique.
(b) What two properties must a decision problem possess in order for it to be NP-complete?
(c) An instance of the Quadratic Residue (QR) decision problem is a triple ( $a, c, m$ ) of positive integers, where $a, c \leq m$, and the problem is to decide if there is an $1 \leq x \leq c$ for which $x^{2} \equiv a \bmod m$. For a given instance $(a, c, m)$ of QR , describe a certificate in relation to $(a, c, m)$.

LO5. Provide the instructions for a URM program that computes the predicate function $f(x, y)=$ $x-y$, where $x-y=0$ in case $y>x$. For each register that is used by your program, described its purpose in one or more sentences.

LO6. Solve the following problems.
(a) Provide the URM program $P$ whose Gödel number equals

$$
2^{30}+2^{74}+2^{112}+2^{137}-1
$$

Show all work.
(b) Let $P$ be a URM program that accepts two inputs $x$ and $y$. What does it mean for $P$ to be universal?
(c) A universal program $P_{U}$ is simulating a program that has 113 instructions and whose Gödel number is

$$
x=2^{3}+2^{179}+2^{191}+2^{196}+2^{224}+2^{268}+\cdots+2^{c_{113}}-1 .
$$

If the current configuration of the computation of $P_{x}$ on some input has encoding

$$
\sigma=2^{5}+2^{14}+2^{16}+2^{23}+2^{28}-1
$$

then provide the next configuration of the computation and its encoding.

LO7. Answer the following. Note: correctly answering two of the three constitutes a pass.
(a) Describe what it means to be a positive instance of the Self Accept decision problem.
(b) In applying the diagonalization method towards proving the undecidability of Self Accept, we defined a computable function $g(x)$ in terms of $f(x)$, the decision function for Self Accept, which we assume to be total computable. Provide the formula for computing $g(x)$ and describe what needs to be established about $g(x)$ in order for the diagonalization method to be successful.
(c) Suppose URM program $P_{86}$ computes $\phi_{86}(x)$ and that

$$
\phi_{86}(x)= \begin{cases}\lfloor x / 5\rfloor & \text { if } x \text { is divisible by } 5 \\ \uparrow & \text { otherwise }\end{cases}
$$

What is the value of $g(86)$ ? Show that $g \neq \phi_{86}$.
LO8. An instance of the decision problem ONTO is a Gödel number $x$, and the problem is to decide if function $\phi_{x}$ is an onto function, meaning that its range is equal to $\mathcal{N}$. Consider the function

$$
g(x)= \begin{cases}1 & \text { if } \phi_{x} \text { is onto } \\ 0 & \text { otherwise }\end{cases}
$$

(a) Evaluate $g(x)$ for each of the following Gödel number's $x$. Note: 2 out of 3 correct is considered passing. Justify your answers.
i. $x=e_{1}$, where $e_{1}$ is the Gödel number of the program that computes the function $\phi_{e_{1}}(y)=2 y$.
ii. $x=e_{2}$, where $e_{2}$ is the Gödel number of the program that computes the function $\phi_{e_{2}}(y)=2 y^{2}$.
iii. $x=e_{3}$, where $e_{3}$ is the Gödel number of the program that computes the function $\phi_{e_{3}}(y)=y$.
(b) Prove that $g(x)$ is not URM computable. In other words, there is no URM program that, on input $x$, always halts and either outputs 1 or 0 , depending on whether or not $\phi_{x}$ onto. Do this by writing a program $P$ that uses $g$ and makes use of the self programming concept. Then explain how $P$ creates a contradiction.

LO9. Solve the following.
(a) Provide the state diagram for a DFA $M$ that accepts all words that contain exactly two 0 's and at least two 1 's.
(b) Show the computation of $M$ on inputs i) $w_{1}=010111$ and ii) $w_{2}=01010$.

LO10. Do the following for the NFA $N$ whose state diagram is shown below.

(a) Provide a table that represents $N$ 's $\delta$ transition function.
(b) Use the table from part a to convert $N$ to an equivalent DFA $M$ using the method of subset states. Draw M's state diagram.
(c) Show the computation of $M$ on input $w=11001$.

LO11. Provide a regular expression that represents the language of all words that contain exactly two 0 's and at least two 1's.

