

Rules for Completing the Problems

NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION allowed when solving these problems. Make sure all these items are put away BEFORE looking at the problems. FAILURE TO ABIDE BY THESE RULES MAY RESULT IN A FINAL COURSE GRADE OF F.

Directions

Choose up to **six problems** to solve. Clearly mark which problems you want us to grade by placing X's in the middle "Grade?" row. **If you don't mark any problems for us to grade or mark 7 or more problems, then we will record grades for the 6 that received the *fewest* points.**

Problem	1	2	3	4	5	6
Grade?						
Points/P/NP						

Problem	LO1	LO2	LO3	LO4	LO5	LO6	LO7	LO8	LO9
Grade?									
P/NP									

3. Answer the following.

a. Provide the definition of what it means for decision problem A to be *Turing reducible* to decision problem B . (10 pts)

b. Suppose $A \leq_T B$ via the logical formula

$$x \in A \Leftrightarrow (\text{query}_B(x^2) \wedge \text{query}_B(x^2 + 3x)) \vee (x \geq 100),$$

and $B \leq_T C$ via the logical formula

$$x \in B \Leftrightarrow (\text{query}_C(3x + 2) \vee (x \leq 75)) \wedge \text{query}_C(5x^3).$$

Prove a logical formula that establishes the fact that $A \leq_T C$. Hint: e.g. $\text{query}_B(x^2)$ is a query to the B -oracle regarding whether or not $x^2 \in B$. (15 pts)

4. Answer the following. Note: correctly solving at least three parts counts for passing LO11.

- a. Recall the mapping reduction $f(\mathcal{C}) = (G, k)$, where f maps an instance of 3SAT to an instance of the Clique decision problem. Given 3SAT instance

$$\mathcal{C} = \{(x_1, x_2, x_4), (\bar{x}_2, x_3, \bar{x}_4), (\bar{x}_1, x_2, x_3), (\bar{x}_1, \bar{x}_3, \bar{x}_4)\},$$

how many vertices does G have? Explain. (5 pts)

- b. Continuing from part a, what is the value of k ? Does G have a k -clique? Defend your answer and provide a k -clique if one exists. (10 pts)

- c. Recall the mapping reduction $f(\mathcal{C}) = (G, a, b)$, where f maps an instance of 3SAT to an instance of the DHP decision problem. Given 3SAT instance

$$\mathcal{C} = \{c_1 = (x_1, x_2, x_4), c_2 = (\bar{x}_2, x_3, \bar{x}_4), c_3 = (\bar{x}_1, x_2, x_3), c_4 = (\bar{x}_1, \bar{x}_3, \bar{x}_4)\}$$

how many “island” vertices does G have? Explain. (5 pts)

- d. Continuing from part c, consider the three vertices lc_3 , rc_3 , and c_3 , where lc_3 and rc_3 belong to the x_2 -diamond. Show how these three vertices are connected to each other. Hint: a variable being set to 0 is analogous to moving from left to right through the diamond. (10 pts)

5. Circle the correct answer (8 pts each) for each multiple choice question. Note: correctly answering at least two counts for passing LO12.

- a. Which of the following decision problems are most likely in class P?
- An instance of **Triangle** is a simple graph G and the problem is to decide if G has a 3-clique.
 - 3SAT**.
 - DHP**.
 - None of the above are likely in class P.
- b. Which of the following certificates is best suited for checking if some instance G of the **Hamilton Cycle** decision problem is a positive instance.
- Certificate C is an ordering of the vertices of G . For example, if G 's vertices are $\{a, b, c, d\}$, then $dbca$ is one possible certificate.
 - Certificate C is a subset of k vertices of G for some $k \geq 0$.
 - Certificate C is an ordering (defined above) of the edges of G .
 - Certificate C is a subset of k edges of G for some $k \geq 0$.
- c. If we know that $A \leq_m^p B$ and A is *not* in P, then it must be true that
- A is NP-complete.
 - B is not in P.
 - B is in NP.
 - None of the above are necessarily true.

6. Given a simple graph $G = (V, E)$, a **dominating set** for G is a set of vertices $D \subseteq V$ for which every vertex $v \in V$ is either a member of D , or is adjacent to some vertex in D .
- If C is a vertex cover for G , explain why it is also a dominating set for G . (5 pts)
 - Give an example of a graph G that has a dominating set D , but D is *not* a vertex cover. Defend your example. (5 pts)
 - Describe a mapping reduction $f : \text{Vertex Cover} \rightarrow \text{Dominated Set}$ from the **Vertex Cover** decision problem to the **Dominated Set** decision problem. For example, an instance of **Dominated Set** is a pair (G, k) , where G is a simple graph and $k \geq 0$ is an integer, and the problem is to decide if G has a dominating set of size k . Clearly define $f(G, k)$ and provide an example of your reduction. Hint: make triangles. (15 pts)