

**CECS 329, Learning Outcome Assessment 9, April 13th, Spring 2023,
Dr. Ebert**

NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION ALLOWED. Submit solutions to at most 2 LO problems on separate sheets of paper.

Problems

LO5. Let x be a natural number, and let $\text{reverse}(x)$ denote the number whose binary representation is that of x 's written backwards. For example, $(12)_2 = 1100$ and so $\text{reverse}(12) = 3$, since

$$(1100)^r = 0011 = (2 + 1)_2 = (3)_2.$$

As another example, $(26)_2 = 11010$ and so $\text{reverse}(x) = 11$, since $(11)_2 = 01011$. Show that $\text{reverse}(x)$ is primitive recursive. You may use any PR functions from the General Models of Computation lecture examples and exercises, including the function $\text{bit}(x, i)$ which equals the i th bit of x when x is written as a binary number, and $\text{len}(x)$ which gives the length of x when written as a binary number. Hint: you do *not* have to provide a recursive definition.

LO6. Do the following.

- (a) Compute the Gödel number for program $P = J(1, 2, 3), S(5), Z(3), T(5, 2)$. Write your answer as a sum of powers of two minus 1.
- (b) Provide the URM program P whose Gödel number equals

$$2^{39} + 2^{230} + 2^{240} + 2^{269} - 1.$$

LO7. Answer and solve the following.

- (a) When simulating a program P_x on input y via a universal program P_U , it is essential that P_U have access to each of the following values during the entire simulation of $P_x(y)$, *except* for
 - i. input y .
 - ii. the number of instructions s of P_x .
 - iii. the index pci for where to locate the program counter in a configuration of $P_x(y)$.
 - iv. All of the above are essential values that P_U must access during the entire computation.
- (b) A universal program P_U is simulating a program that has 137 instructions and whose Gödel number is

$$x = 2^8 + 2^{22} + 2^{726} + 2^{751} + 2^{c_5} + \dots + 2^{c_{137}} - 1.$$

If the current configuration of the computation of P_x on some input has encoding

$$\sigma = 2^5 + 2^7 + 2^{13} + 2^{16} + 2^{20} - 1,$$

then provide the next configuration of the computation *and* its τ encoding. Hint: $176 = 16 \times 11$.

LO8. Do the following.

- (a) In one or more complete sentences, describe what is asserted by the Church-Turing Thesis.
- (b) Consider the function

$$g(x) = \begin{cases} 1 & \text{if } \phi_x \text{ is total} \\ 0 & \text{otherwise} \end{cases}$$

In other words, $g(x) = 1$ iff program P_x halts on all of its inputs. We want to prove that $g(x)$ is undecidable, meaning there is no URM program that computes g . To do this, let's assume that $g(x)$ is computable by some URM program G . Then define the function $f(x)$ as follows.

$$f(x) = \begin{cases} P_x(x) + 1 & \text{if } g(x) = 1 \\ 0 & \text{if } g(x) = 0 \end{cases}$$

Use part a and an informal description of how you would go about computing $f(x)$ to establish that $f(x)$ is URM computable.

- (c) Let e denote the Gödel number of the URM program F that computes $f(x)$ from part b. In other words $F = P_e$. Show that a contradiction arises when we try to compute $F(e) = P_e(e)$.

LO9. An instance of the decision problem **Zero** is a Gödel number x , and the problem is to decide if function ϕ_x equals the zero function, i.e. the function that outputs 0 on every input. Consider the function

$$g(x) = \begin{cases} 1 & \text{if } \phi_x(y) = 0 \text{ for all inputs } y \\ 0 & \text{otherwise} \end{cases}$$

- (a) Evaluate $g(x)$ for each of the following Gödel number's x . Note: 2 out of 3 correct is considered passing.
 - i. $x = e_1$, where e_1 is the Gödel number of the program $P = T(1, 2), Z(1), S(2)$.
 - ii. $x = e_2$, where e_2 is the Gödel number of the program $P = T(1, 2), Z(1), Z(2), J(1, 2, 1)$.
 - iii. $x = e_3$, where e_3 is the Gödel number of the program that computes $g(x)$.

Hint: $T(m, n)$ means copy the value of R_m to R_n .

- (b) Prove that $g(x)$ is not URM computable. In other words, there is no URM program that, on input x , always halts and either outputs 1 or 0, depending on whether or not ϕ_x equals the zero function. Do this by writing a program P that uses g and makes use of the **self** programming concept. Then show how P creates a contradiction.