

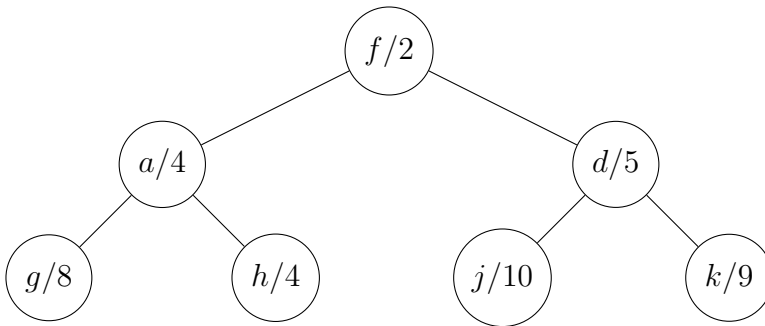
**NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION ALLOWED.** Submit each solution on a separate sheet of paper.

## Problems

LO6. The tree below shows the state of the binary min-heap at the beginning of some round of Dijkstra’s algorithm, applied to some weighted graph  $G$ . If  $G$  has directed edges

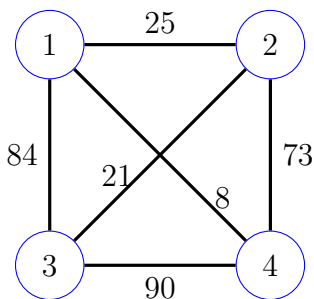
$$(a, f, 1), (d, g, 2), (f, g, 5), (f, h, 2), (f, k, 3),$$

then draw a plausible state of the heap at the end of the round.



LO7. Do the following. Note: correctly solving this problem counts for passing LO7.

- (a) The dynamic-programming algorithm that solves the **Runaway Traveling Salesperson** optimization problem (Exercise 30 from the Dynamic Programming Lecture) defines a recurrence for the function  $mc(i, A)$ . In words, what does  $mc(i, A)$  equal? Hint: do *not* write the recurrence (see Part b). Hint: we call it “Runaway TSP” because the salesperson does *not* return home.
- (b) Provide the dynamic-programming recurrence for  $mc(i, A)$ .
- (c) Apply the recurrence from Part b to the graph below in order to calculate  $mc(1, \{2, 3, 4\})$  Show all the necessary computations and use the solutions to compute an optimal path for the salesperson.



LO8. Do/answer the following.

- (a) Draw the implication graph  $G_C$  associated with the 2SAT instance

$$\mathcal{C} = \{(\bar{x}_1, x_3), (x_2, x_3), (x_2, \bar{x}_3), (\bar{x}_2, \bar{x}_3), (x_3, \bar{x}_5), (x_4, x_5)\}.$$

- (b) Apply the improved 2SAT algorithm to obtain a satisfying assignment for  $\mathcal{C}$ . When deciding on the next reachability set  $R_l$  to compute, follow the literal order  $l = x_1, \bar{x}_1, \dots, x_5, \bar{x}_5$ . For each consistent reachability set encountered, provide the partial assignment  $\alpha_{R_l}$  associated with  $R_l$  and draw the reduced implication graph before continuing to the next reachability set. Note: do *not* compute the reachability set for a literal that has already been assigned a truth value. Provide a final assignment  $\alpha$  and verify that it satisfies all the clauses.
- (c) Suppose 2SAT instance  $\mathcal{C}$  is satisfiable and uses 115 variables and 336 clauses. Using the original 2SAT algorithm, what is the *least* number of queries to a Reachability oracle that needs to be made in order to establish  $\mathcal{C}$ 's satisfiability. In other words, if we make fewer than this number of queries then it is possible that the 2SAT instance  $\mathcal{C}$  may be unsatisfiable. Explain.

LO9. Answer the following questions?

- (a) Provide the definition of what it means to be a mapping reduction from problem  $A$  to problem  $B$ .
- (b) As part of her network security project, Julie is working with a simple graph  $H$  that has 145 vertices and 372 edges. She needs to know whether or not  $H$  has a Hamilton path. Her colleague Simon has implemented the Python function

```
Boolean has_LPath(Graph G, int k);
```

that decides if the input graph  $G$  has a simple path of length  $k$ . Explain how Julie can use Simon's function to answer her question. If used properly, why will Simon's function return the answer that Julie seeks.

LO10. An instance of the **Feedback Arc Set (FAS)** decision problem is a pair  $(G, k)$  where  $G = (V, E)$  is a directed graph and  $k$  is a nonnegative integer. The problem is to decide if there is a subset  $E'$  of  $k$  edges that can be removed from  $G$  so that  $G - E'$  is an acyclic graph.

- (a) For a given instance  $(G, k)$  of FAS, describe a certificate in relation to  $(G, k)$ .
- (b) Provide a semi-formal verifier algorithm that takes as input i) an instance  $(G, k)$ , ii) a certificate for  $(G, k)$  as defined in part a, and decides if the certificate is valid for  $(G, k)$ . Do this by making use of a reachability oracle via function `reachable( $H, a, b$ )` which returns true iff vertex  $b$  is reachable from vertex  $a$  in graph  $H$ . Hint: for each vertex  $v \in V$ , consider the set  $N^+(v) = \{u \mid (u, v) \in E\}$ .
- (c) Provide size parameters that may be used to measure the size of an instance  $(G, k)$  of FAS.
- (d) Use the size parameters from part c to describe the running time of your verifier from part b. Assume that a single call to function `reachable` requires a linear number of steps with respect to the size of the input graph. Defend your answer in relation to the algorithm you provided for the verifier.