## CECS 528, Learning Outcome Assessment 10, Pink, Fall 2023, Dr. Ebert

## NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION ALLOWED. Submit each solution on a separate sheet of paper.

## Problems

LO6. Recall the use of the disjoint-set data structure for the purpose of improving the running time of the Unit Task Scheduling algorithm. For the set of tasks

| Task | a | b | c | d | e | f |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| Deadline Index | 2 | 1 | 2 | 4 | 3 | 4 |
| Profit | 60 | 50 | 40 | 30 | 20 | 10 |

For each task, show the M-Tree forest after it has been inserted (or at least has attempted to be inserted in case the scheduling array is full). Note that the earliest possible deadline index is 0 , meaning that the earliest slot in the schedule array has index 0 . Also, assume that an insert attempt that takes place at index $i$ results in the function call root $(i)$. Finally, to receive credit, your solution should show six different snapshots of the M-Tree forest.

LO7. Answer the following.
(a) The dynamic-programming algorithm that solves the 0-1 Knapsack optimization problem defines a recurrence for the function $p(i, c)$. In words, what does $p(i, c)$ equal? Hint: do not write the recurrence (see Part b).
(b) Provide the dynamic-programming recurrence for $p(i, c)$.
(c) Apply the recurrence from Part b to a knapsack having capacity $M=10$ and items

| item | weight | profit |
| :--- | :--- | :--- |
| 1 | 5 | 30 |
| 2 | 4 | 30 |
| 3 | 1 | 20 |
| 4 | 4 | 40 |
| 5 | 5 | 30 |
| 6 | 5 | 60 |

Show the matrix of subproblem solutions and use it to provide an optimal set of items.
LO8. Do/answer the following.
(a) Draw the implication graph $G_{\mathcal{C}}$ associated with the 2SAT instance

$$
\mathcal{C}=\left\{\left(x_{1}, \bar{x}_{5}\right),\left(x_{1}, x_{5}\right),\left(\bar{x}_{2}, \bar{x}_{3}\right),\left(\bar{x}_{2}, \bar{x}_{4}\right),\left(x_{2}, x_{4}\right),\left(x_{2}, \bar{x}_{4}\right),\left(\bar{x}_{3}, x_{4}\right),\left(x_{3}, x_{5}\right)\right\} .
$$

(b) Apply the improved 2SAT algorithm to obtain a satisfying assignment for $\mathcal{C}$. When deciding on the next reachability set $R_{l}$ to compute, follow the literal order $l=x_{1}, \bar{x}_{1}, \ldots, x_{5}, \bar{x}_{5}$. For
each consistent reachability set encountered, provide the partial assignment $\alpha_{R_{l}}$ associated with $R_{l}$ and draw the reduced implication graph before continuing to the next reachability set. Note: do not compute the reachability set for a literal that has already been assigned a truth value. Provide a final assignment $\alpha$ and verify that it satisfies all the clauses.
(c) Suppose 2SAT instance $\mathcal{C}$ is unsatisfiable and has 234 variables and 754 clauses. Using the original 2SAT algorithm, what is the least number of queries to a Reachability oracle that needs to be made in order to establish $\mathcal{C}$ 's unsatisfiability. In other words, it is possible that $\mathcal{C}$ could be proved unsatisfiable after this number of queries.

LO9. Answer the following.
(a) Provide the definition of what it means to be a mapping reduction fromm problem $A$ to problem $B$.
(b) Suppose $(G, k=3)$ is an instance of the Vertex Cover decision problem, where $G$ is drawn below. Draw $f(G, k)$, where $f$ is the mapping reduction from Vertex Cover to the Half Vertex Cover decision problem.

(c) Verify that $f$ is valid for input $(G, k)$ in the sense that both $(G, k)$ and $f(G)$ are either both positive instances or both negative instances of their respective problems. Defend your answer. Hint: it takes two vertices to cover each edge of a triangle in a graph.

LO10. An instance of the Boolean Vector Sum (BVS) decision problem is a pair $(B, k)$ where $B$ is a set of $n$ Boolean vectors, each having length equal to $m>0$, and $k$ is a nonnegative integer. The problem is to decide if $B$ has a subset of $k$ vectors $\left\{v_{1}, \ldots, v_{k}\right\}$ for which

$$
v_{1} \vee v_{2} \vee \cdots \vee v_{k}=(\underbrace{1, \ldots, 1}_{m \text { times }}),
$$

where operation $\vee$ represents bitwise OR.
(a) For a given instance $(B, k)$ of BVS describe a certificate in relation to $(B, k)$.
(b) Provide a semi-formal verifier algorithm that takes as input i) an instance ( $B, k$ ), ii) a certificate for $(B, k)$ as defined in part a, and decides if the certificate is valid for $(B, k)$.
(c) Provide size parameters that may be used to measure the size of an instance of BVS.
(d) Use the size parameters from part c to describe the running time of your verifier from part b. Defend your answer in relation to the algorithm you provided for the verifier.

