

## Rules for Completing the Problems

**NO NOTES, BOOKS, ELECTRONIC DEVICES, OR INTERPERSONAL COMMUNICATION** allowed when solving these problems. Make sure all these items are put away **BEFORE** looking at the problems. **FAILURE TO ABIDE BY THESE RULES MAY RESULT IN A FINAL COURSE GRADE OF F.**

## Directions

Choose up to **six problems** to solve. Clearly mark each problem you want graded by placing an 'X' or check mark in the appropriate box. **If you don't mark any problems for us to grade or mark 7 or more problems, then we will record grades for the 6 that received the *fewest* points.**

Problem	1	2	3	4	5	6
Grade?						
Result						

Your Full Name:

Your Class ID:

1. Answer the following. Note: correctly solving this problem counts for passing LO1.

a. Provide the definition of what it means to be a mapping reduction from decision problem  $A$  to decision problem  $B$ . (5 pts)

b. For the mapping reduction  $f : \text{Subset Sum} \rightarrow \text{Set Partition}$ , based on the reduction provided in lecture, determine  $f(S, t)$  for **Subset Sum** instance

$$(S = \{4, 7, 15, 19, 22, 38, 44, 45\}, t = 111).$$

Show work. (10 pts)

c. Verify that  $A = \{22, 44, 45\}$  is a valid certificate for  $(S, t)$ , and use  $A$  to construct a valid certificate for  $f(S, t)$ . Explain and show work. (10 pts)

2. An instance of **Set Cover** is a triple  $(\mathcal{S}, m, k)$ , where  $\mathcal{S} = \{S_1, \dots, S_n\}$  is a collection of  $n$  subsets, where  $S_i \subseteq \{1, 2, \dots, m\}$ , for each  $i = 1, \dots, n$ , and  $k$  is a nonnegative integer. The problem is to decide if there are  $k$  subsets  $S_{i_1}, \dots, S_{i_k}$  for which

$$S_{i_1} \cup \dots \cup S_{i_k} = \{1, 2, \dots, m\}.$$

Note: correctly solving this problem counts for passing LO2.

- a. Verify that  $(\mathcal{S}, m, k)$  is a positive instance of **Set Cover**, where  $m = 9$ ,  $k = 4$ , and

$$\mathcal{S} = \{\{1, 3, 5\}, \{3, 7, 9\}, \{2, 4, 5\}, \{2, 6, 7\}, \{6, 7, 9\}, \{2, 7, 9\}, \{1, 3, 7\}, \{4, 5, 8\}\}.$$

(5 pts)

- b. For a given instance  $(\mathcal{S}, m, k)$  of **Set Cover** describe a certificate in relation to  $(\mathcal{S}, m, k)$ .

(5 pts)

- c. Provide a semi-formal verifier algorithm that takes as input i) an instance  $(\mathcal{S}, m, k)$ , ii) a certificate for  $G$  as defined in part b, and decides if the certificate is valid. (5 pts)

- d. Provide size parameters that may be used to measure the size of an instance of **Set Cover**. Hint: there should be two of them. (5 pts)

- e. Use the size parameters from part d to describe the running time of your verifier from part c. Defend your answer in relation to the algorithm you provided for the verifier. (5 pts)

3. Answer the following. Note: correctly solving at least three parts counts for passing LO3.

- a. Recall the mapping reduction  $f(\mathcal{C}) = (S, t)$ , where  $f$  maps an instance of **3SAT** to an instance of the **Subset Sum** decision problem. Given **3SAT** instance

$$\mathcal{C} = \{c_1 = (x_1, x_2, x_4), c_2 = (\bar{x}_2, x_3, \bar{x}_4), c_3 = (\bar{x}_1, x_2, x_3), c_4 = (\bar{x}_1, \bar{x}_3, \bar{x}_4),$$

What is the cardinality of  $S$ ? What is the value of  $t$ ? Explain and/or Show work. (5 pts)

- b. Continuing from part a, is  $(S, t)$  a positive instance? If not, explain why. If yes, provide a subset of  $S$  that sums to target  $t$ . (6 pts)

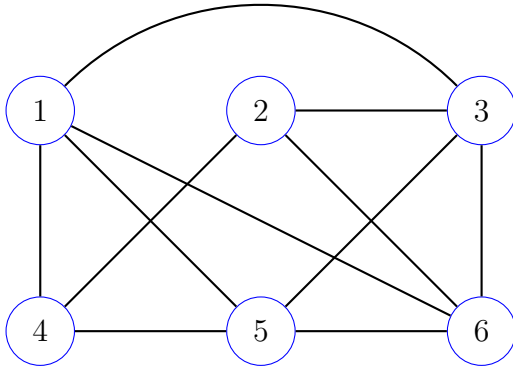
- c. Recall the mapping reduction  $f(G) = (G', k)$ , where  $f$  maps an instance of **Hamilton Cycle** to an instance of the **TSP** decision problem. If  $G$  has 8 vertices, then how many edges does  $G'$  have? What is the value of  $k$ ? Explain. (7 pts)

- d. Continuing from part c, if  $G$  has a Hamilton path, but no Hamilton cycle, then what is the total cost of the least-cost Hamilton cycle in  $G'$ . In other words, determine the cost of the most cost-efficient cycle that appears in  $G'$ . Explain. (7 pts)

4. Answer the following. Note: scoring 14 or more points counts for passing LO4.
- a. Write either P, NP, or co-NP next to each of the following decision problems in terms of which one best characterizes the complexity of the problem. Note: there is only one best answer for each. Two points each.
- i. An instance of **Triangle** is a simple graph  $G$  and the problem is to decide if  $G$  has a 3-clique.
  - ii. An instance of **Tired Traveling Salesperson (TTSP)** is a complete weighted graph  $G$  and an integer  $k \geq 0$ . The problem is to decide if every Hamilton Cycle in  $G$  has a cost that exceeds  $k$ .
  - iii. An instance of **Composite** is a natural number  $n \geq 2$ , and the problem is to decide if  $n$  is a composite number, i.e. is divisible by a number  $2 \leq d \leq n - 1$ .
  - iv. An instance of **Quadratic Diophantine** is a triple of positive integers  $(a, b, c)$  and the problem is to decide if there are positive integers  $x$  and  $y$  for which  $ax^2 + by = c$ .
- b. Circle the correct answer. Which of the following certificates is best suited for checking if some simple graph  $G = (V, E)$  is a positive instance of **Hamilton Cycle**. (8 points)
- i. Certificate  $C$  is a permutation of all  $n$  vertices of  $G$ .
  - ii. Certificate  $C$  is a subset of  $n$  edges of  $G$ , where  $n = |V|$ .
  - iii. Certificate  $C$  is a permutation of all  $m$  edges of  $G$ .
  - iv. Certificate  $C$  is a subset of  $n$  edges of  $G$ , where  $n = |V|$ .
- c. Circle the correct answer. Which of the following is a true statement (9 points)
- i. If  $A \leq_m^p B$ ,  $A \in \text{P}$ , and  $B$  is NP-complete, then  $B \in \text{P}$ .
  - ii. If  $A \leq_m^p B$ ,  $A$  is NP-complete,  $B \in \text{NP}$ , then  $\text{P} \neq \text{NP}$ .
  - iii. If  $A \leq_m^p B$ ,  $A$  is NP-complete,  $B \in \text{P}$ , then  $\text{P} = \text{NP}$ .
  - iv. If  $A \leq_m^p B$ ,  $A \in \text{P}$ , and  $B$  is NP-complete, then  $\text{P} = \text{NP}$ .

5. In lecture we showed a polynomial-time mapping reduction  $f : \text{HP} \rightarrow \text{HC}$ .

- a. Draw the image of  $f(G, 1, 6)$ , where  $G$  is the graph shown below. Hint: you may represent your image by modifying  $G$ . (5 pts)



- b. Provide a complete description of a valid polynomial-time mapping reduction  $g : \text{HC} \rightarrow \text{HP}$  from **Hamilton Cycle** to **Hamilton Path**. Function  $g$  must map a graph  $G$  to a triple  $(G', a, b)$ , where  $a$  and  $b$  are vertices of  $G'$  and  $G$  has a Hamilton Cycle iff there is a Hamilton path from  $a$  to  $b$ . Use the graph from part a to demonstrate your reduction. Hint: clone a vertex of  $G$ . Argue that your reduction is valid. (20 pts)

6. The **Half Clique** decision problem is the problem of deciding if a simple graph  $G = (V, E)$  has a **Clique** of size  $|V|/2$ .

a. Describe an embedding reduction  $f$  from **Half Clique** to **Clique**. (10 points)

b. Now describe a contraction reduction from **Clique** to **Half Clique** and demonstrate it for the **Clique** instance  $(G, k = 4)$ , where  $G$  is the graph shown below. (15 points)

